

# Profile Guided Offline Optimization of Hidden Class Graphs for JavaScript VMs in Embedded Systems

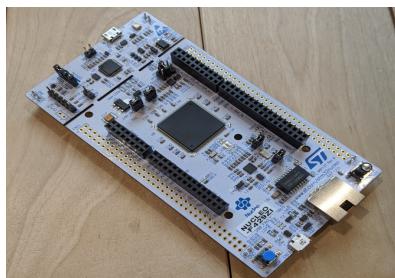
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# JavaScript in IoT

- JavaScript engines for IoT became popular
  - IoT.js, Moddable, **eJSVM**,...
- Challenge: memory footprint
  - Around 256 KB of RAM is available
  - More than 20 KB of RAM is occupied by meta-objects in eJSVM



STM32F429

- Arm Cortex-M4
- **256 KB of SRAM**

## Raspberry Pi pico specification

Dual-core Arm Cortex-M0+ processor, flexible clock

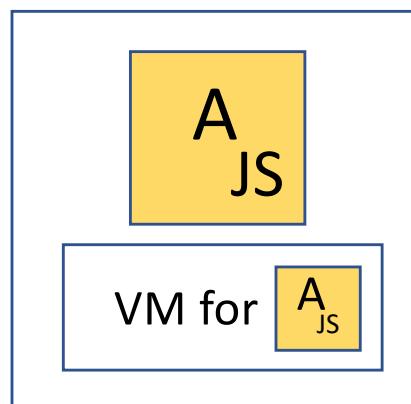
264kB on-chip SRAM

2MB on-board QSPI flash

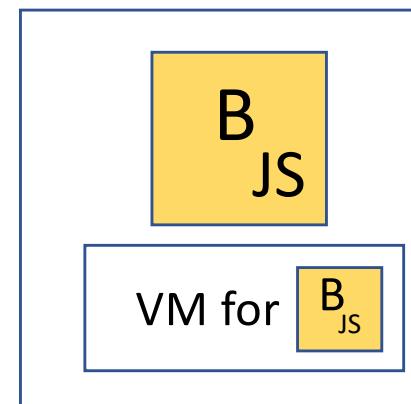
2.4GHz 802.11n wireless LAN (Raspberry Pi Pico W)

# Closed World Assumption

- We can assume program is fixed for a particular IoT product
  - For product A, VM executes only A.js
- VM specialisation to a particular application is feasible



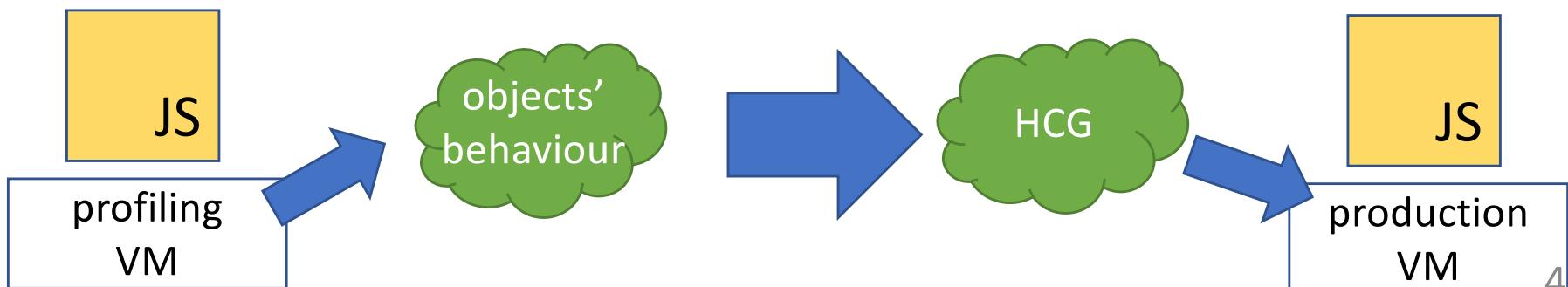
product A



product B

# Overview of Our Work

- Specialise hidden class graph (HCG)
  - HCG represents type information of objects
  - HCG is created and grows during execution in accordance with program's behaviour
- Steps
  1. Collect objects' behaviour from profiling run
  2. Construct a static HCG and optimise it offline
  3. Use static HCG in actual runs



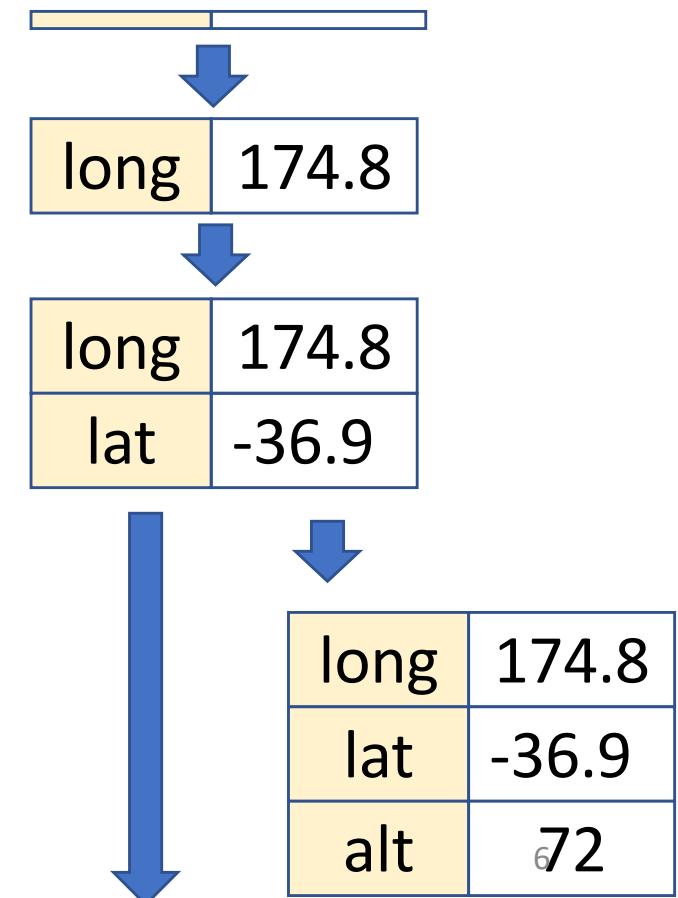
# Agenda

- Introduction
- Hidden classes
- Optimised Hidden Class Construction
- Evaluation

# JavaScript Object

- Not statically typed
  - Properties are added dynamically
  - Set of properties depends on control-flow

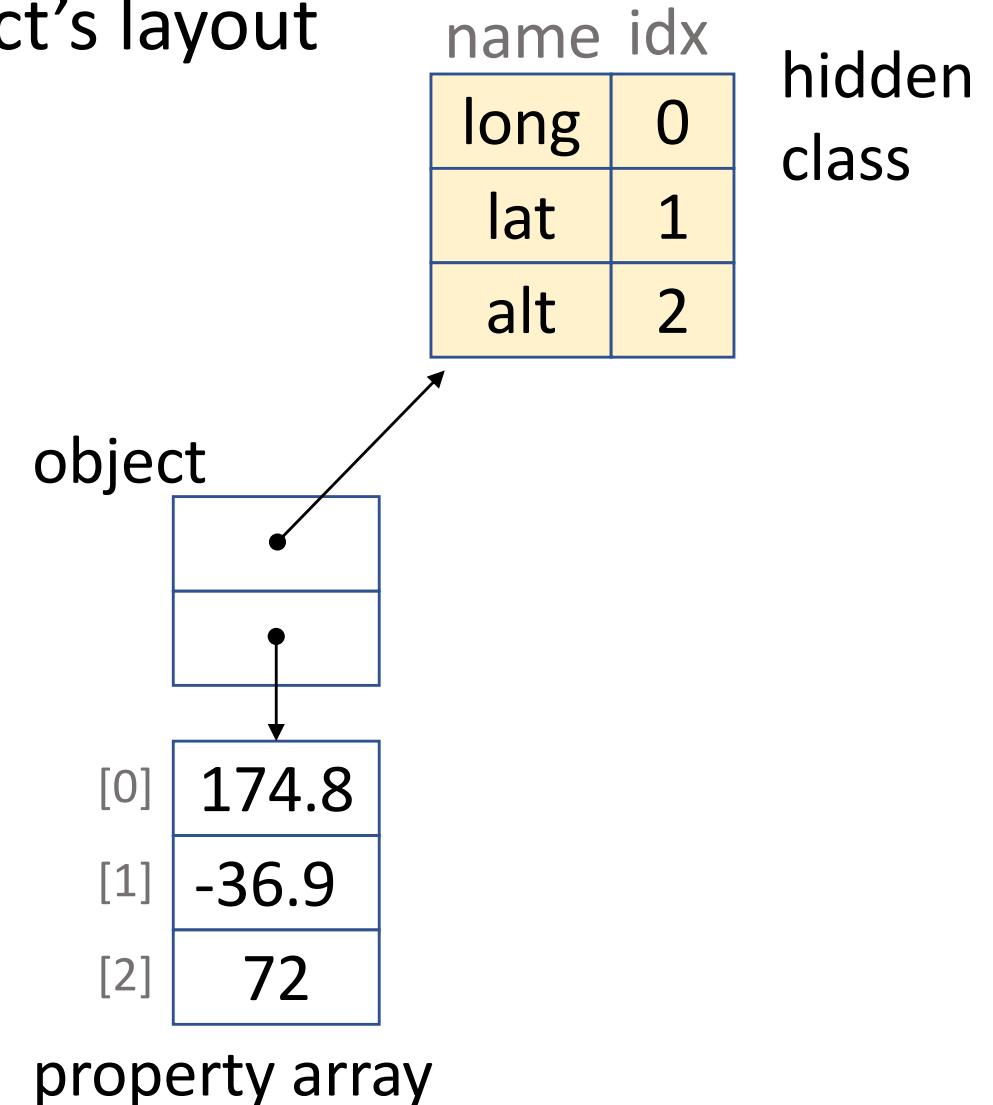
```
readGPS () {  
    let loc = {};  
    loc.long = getLongitude ();  
    loc.lat = getLatitude ();  
    if (hasAltitude ())  
        loc.alt = getAltitude ();  
    return loc;  
}
```



# Hidden Class (HC)

- Meta-object having object's layout
- object = (HC, prop array)

long	174.8
lat	-36.9
alt	72



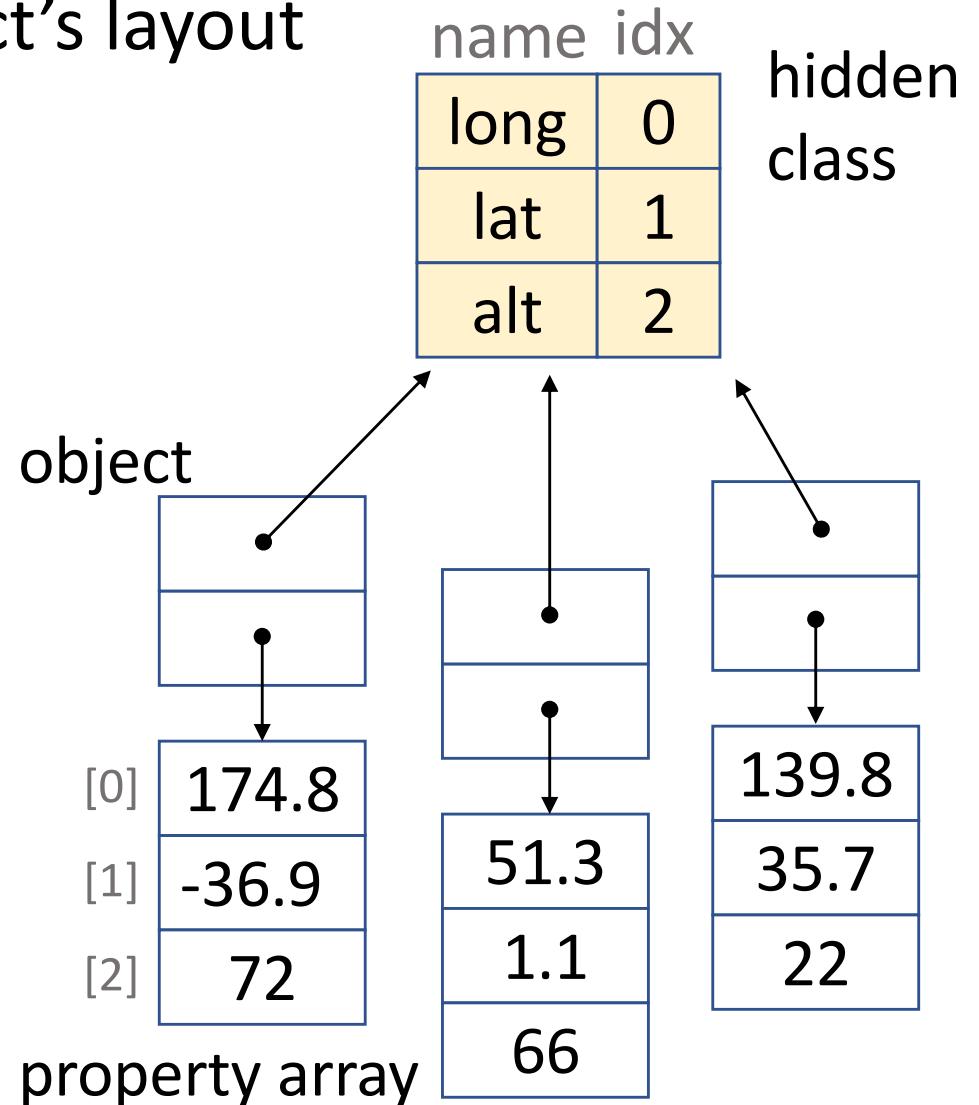
# Hidden Class (HC)

- Meta-object having object's layout
- object = (HC, prop array)
- Shared with all instances

long	174.8
lat	-36.9
alt	72

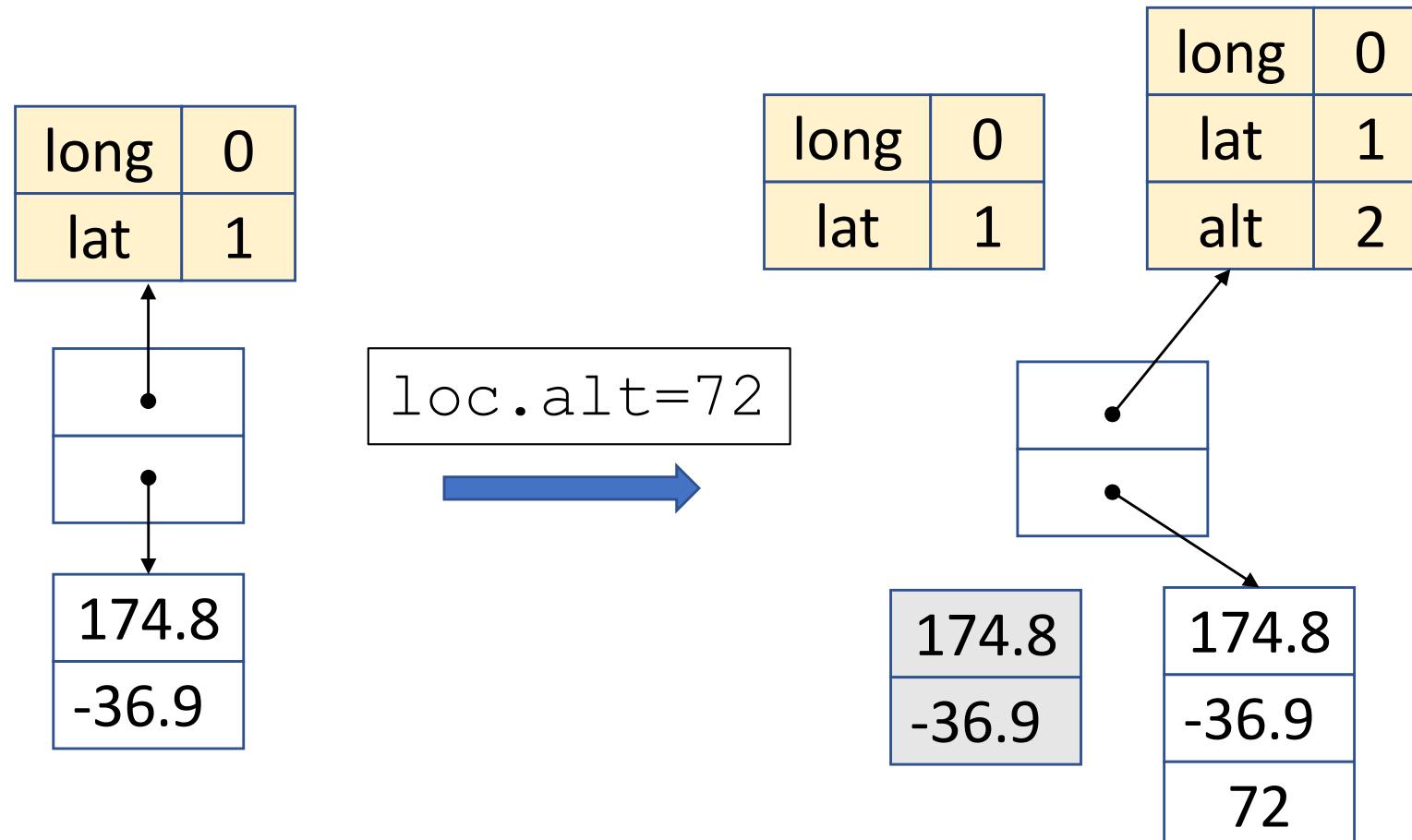
long	174.8
lat	35.7
alt	40

long	51.3
lat	1.1
alt	66



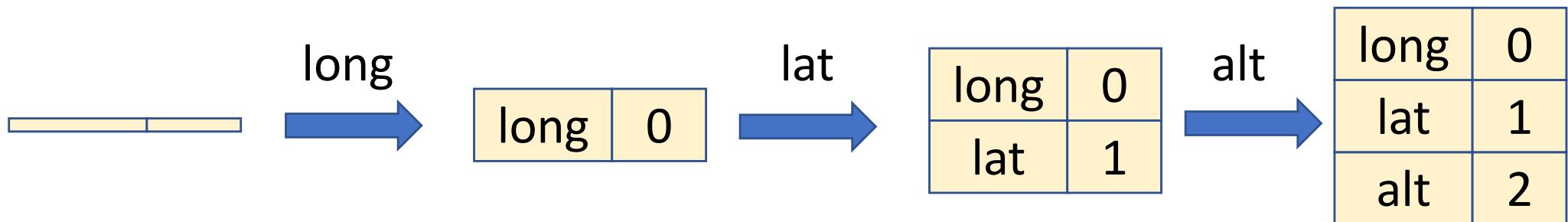
# Hidden Class Transition

- Adding new property causes HC transition
  - Find next HC, or create it if it has not been created
  - Re-allocate property array



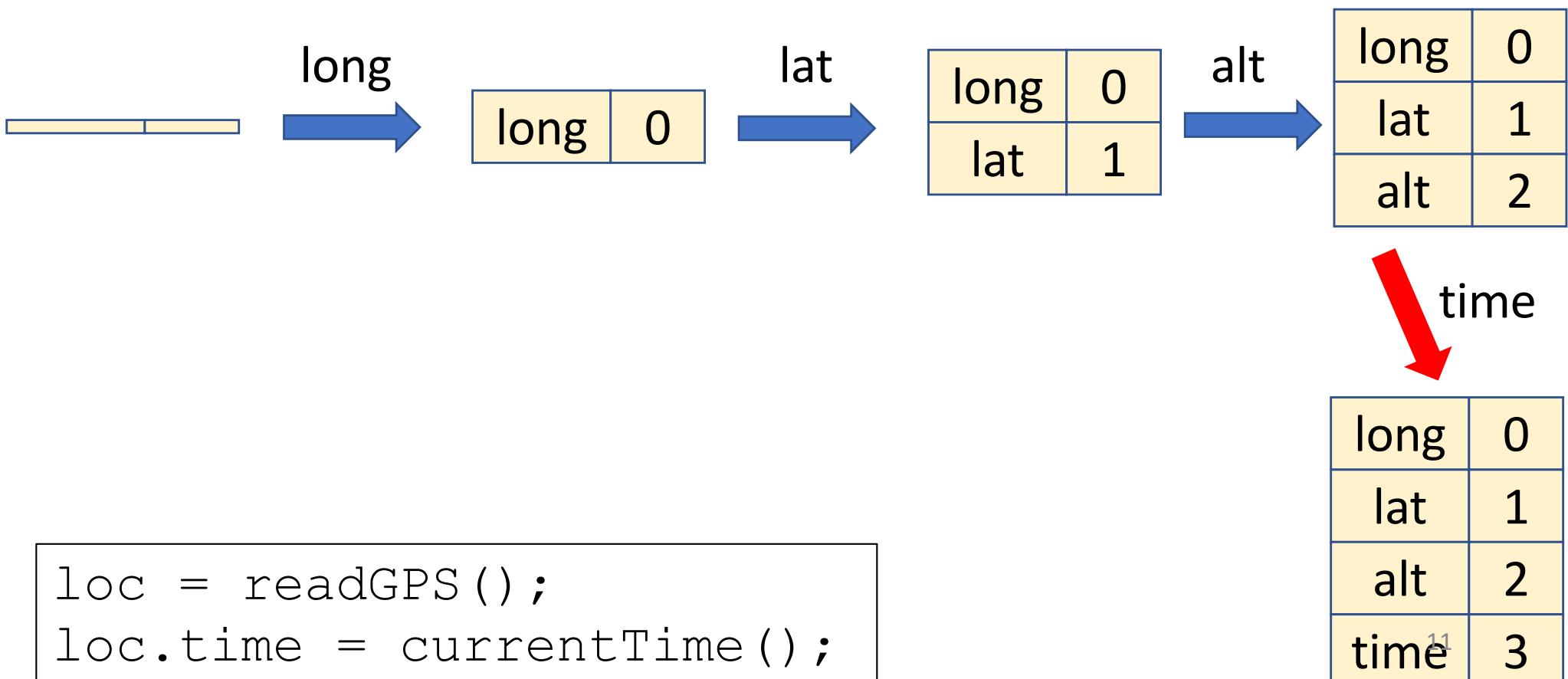
# Hidden Class Graph

- Hidden class graph (HCG) enables to find next HC quickly
  - node: HC
  - edge: transition labelled with property name



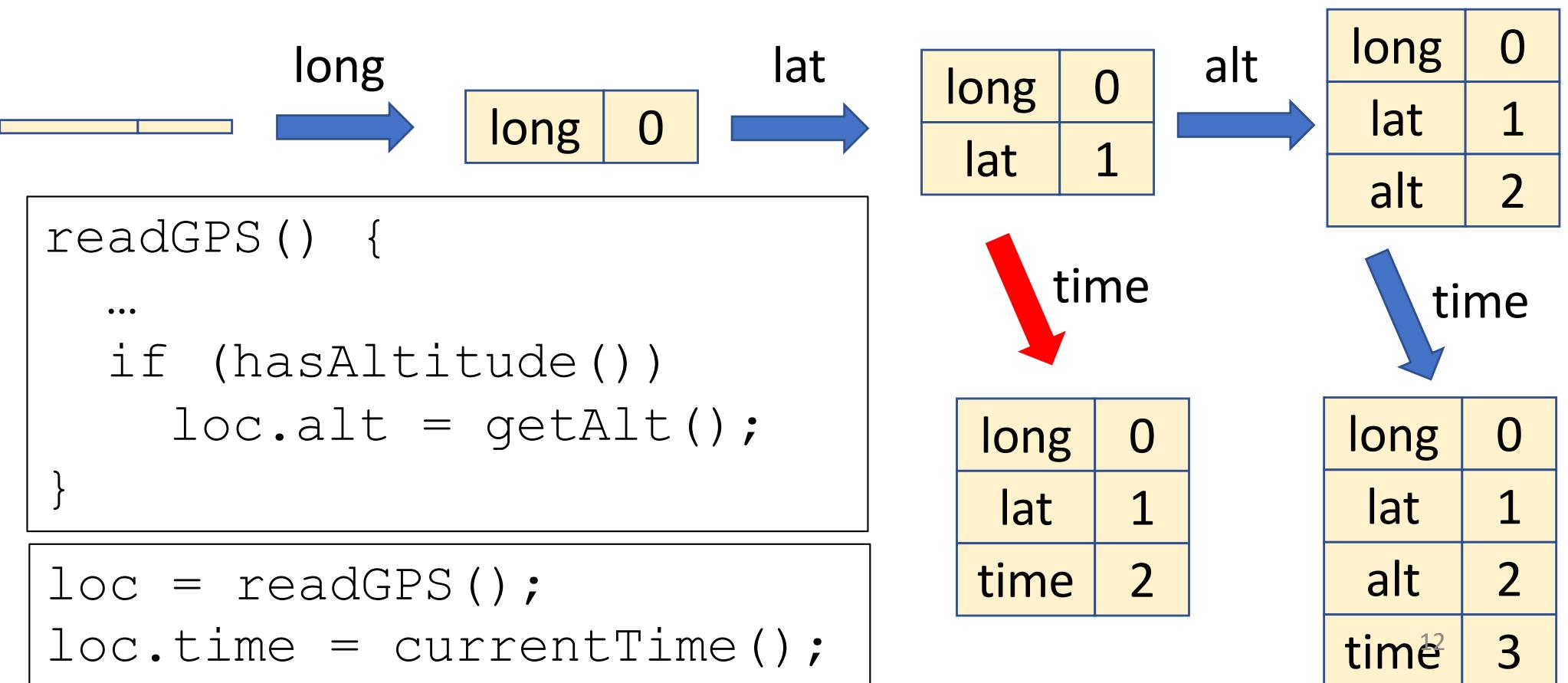
# HCG grows during execution

- New property creates new HC
- New HC is added to HCG



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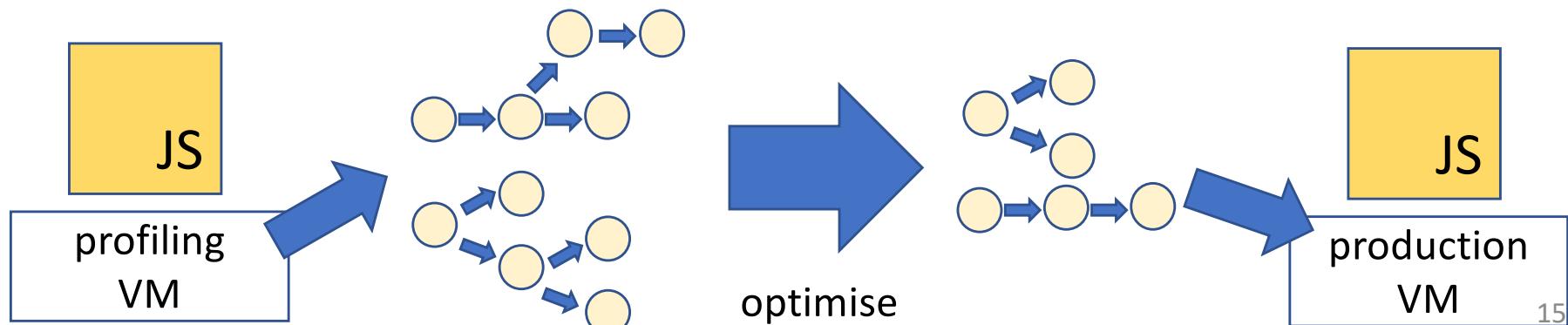


# Agenda

- Introduction
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# Offline Optimisation of HCG

- Optimise HCG with the following policy
  - Reduce memory footprint
    - Shrink HCG and reduce object size
    - Allow small space-inefficiency for speed
  - Use optimised HCG in production VM
    - Run-time optimisation relying on assumption that HCG is stable

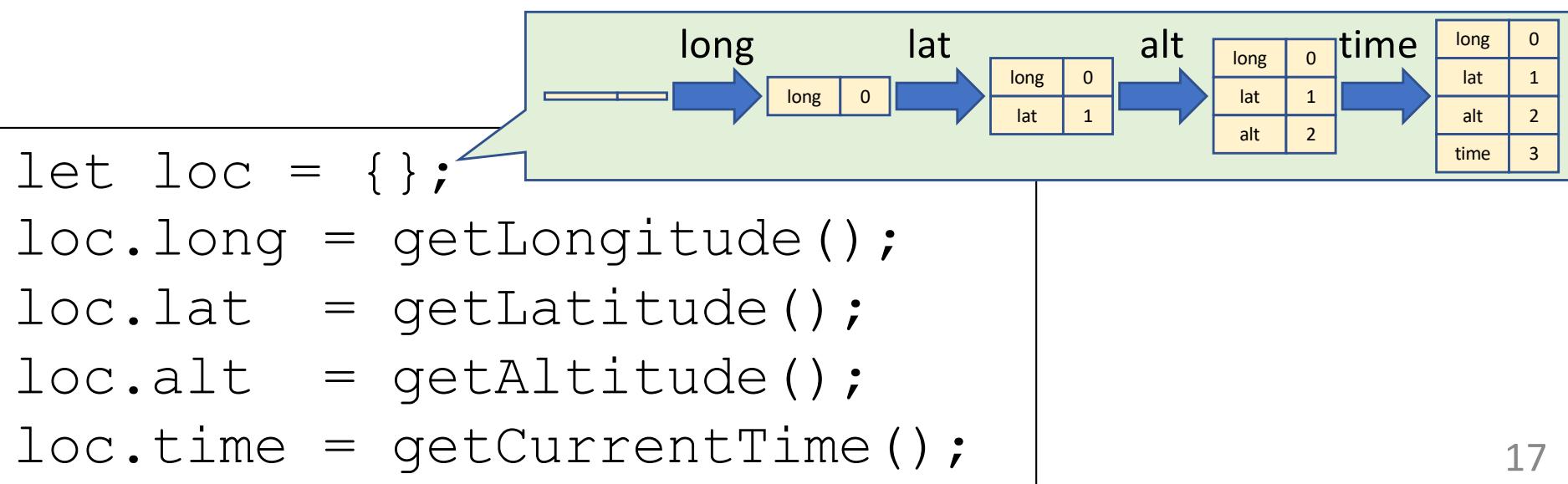


# Optimisations

1. Eliminating intermediate HCs
2. Moving branches

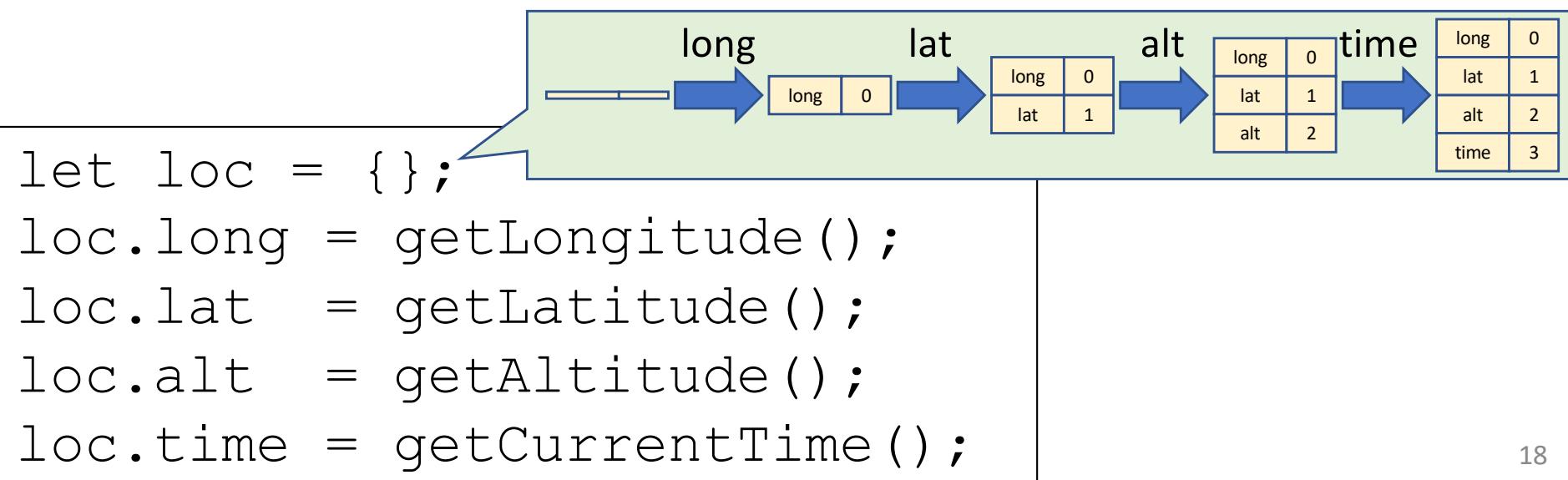
# Layout-monomorphic allocation site

- 95.8 % of allocation sites are layout-monomorphic
- Layout-monomorphic allocation site:  
all objects allocated there obtain the same set of properties in the same order
  - Eventually get transitions to the same HC



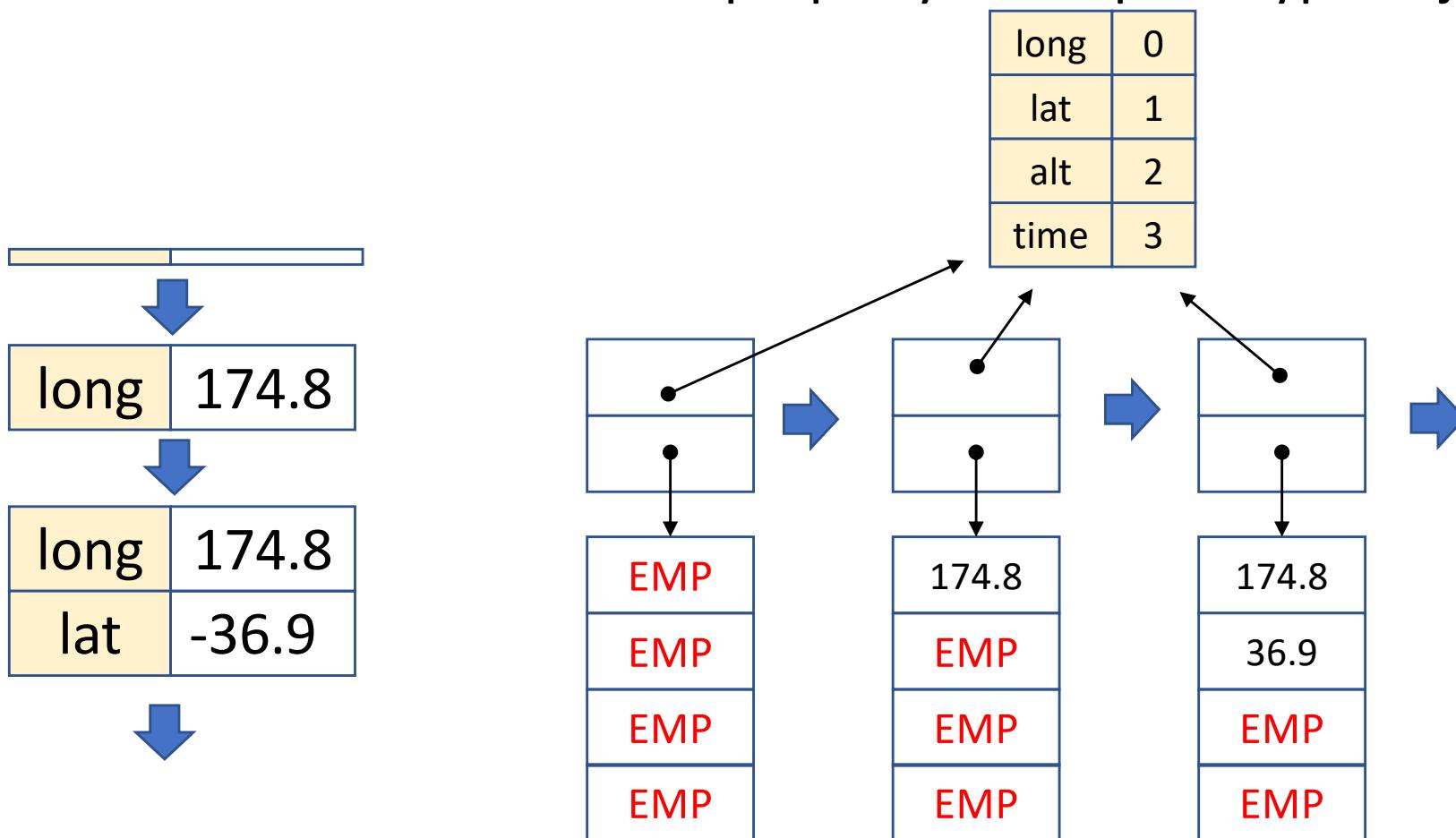
# Optimisation for layout-monomorphism: pre-transitioning

- Eliminate all hidden classes but the last from HCG
- Objects are created with their final layout
  - No re-allocation overhead of property array



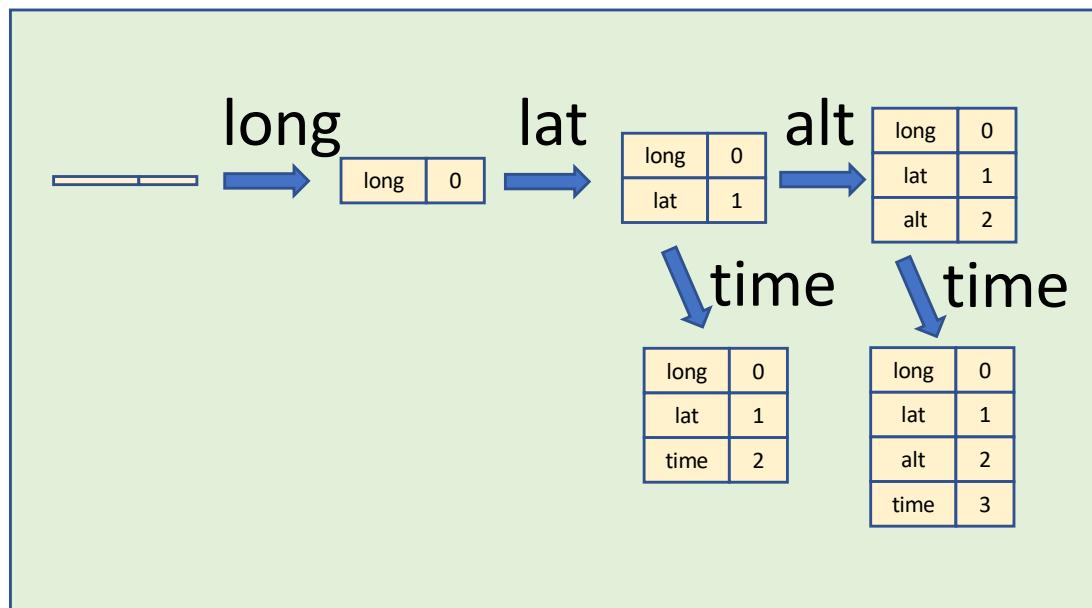
# EMPTY value

- Initialise property array slots with EMP to indicate absence of the property
  - Allow us to search for property in the prototype object.



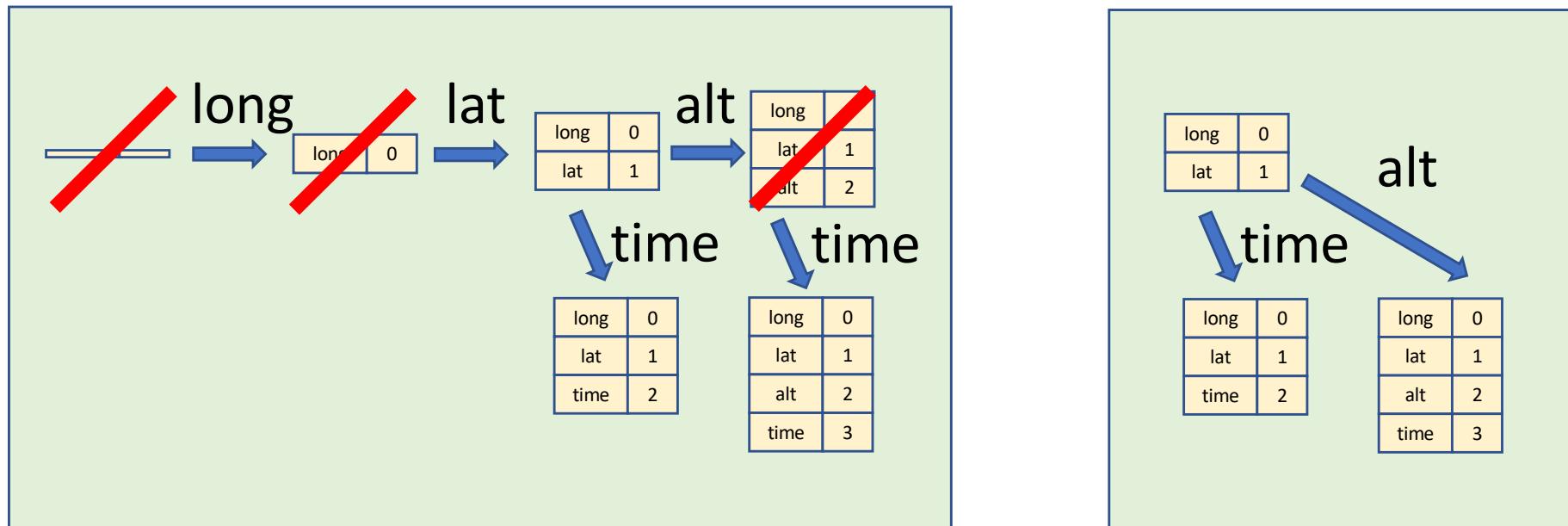
# Optimisation 1: elimination of Intermediate HCs

- Generalization of pre-transitioning
- Eliminate all internal nodes but branching nodes



# Optimisation 1: elimination of Intermediate HCs

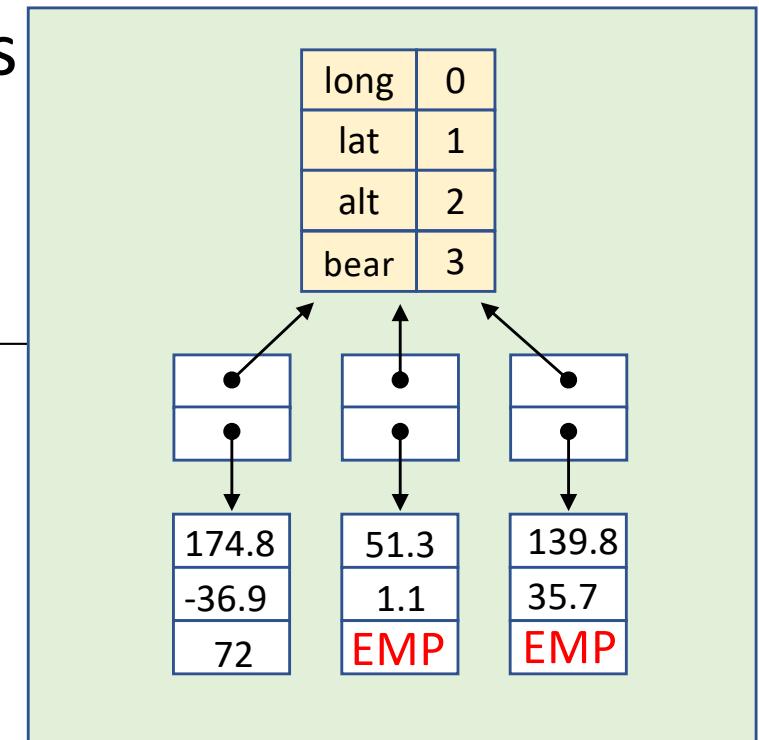
- Generalization of pre-transitioning
- Eliminate all internal nodes but branching nodes



# Over-allocation

- Aggressive elimination increases memory footprint
  - Memory for all possible properties are reserved
- Examples
  - Parts of objects get extra props
  - Props are added in the future

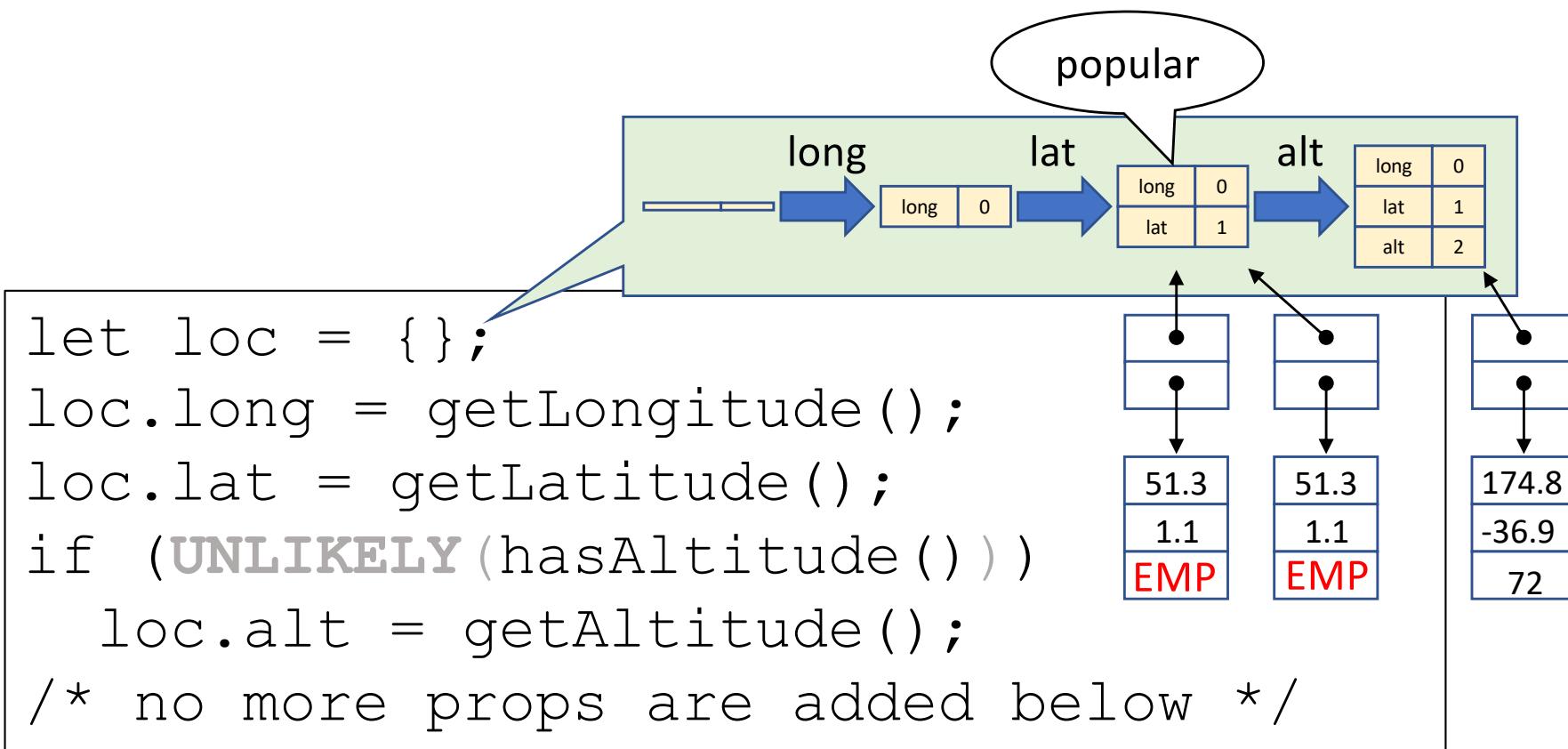
```
let loc = {};  
loc.long = getLongitude();  
loc.lat = getLatitude();  
if (UNLIKELY(hasAltitude()))  
    loc.alt = getAltitude();  
/* no more props are added below */
```



# Optimisation 1': preserve popular HCs

- popular HC:  $\max_{t \in execution} (\text{#of instances}) > K$ 
  - Sample # of instances at each GC cycle

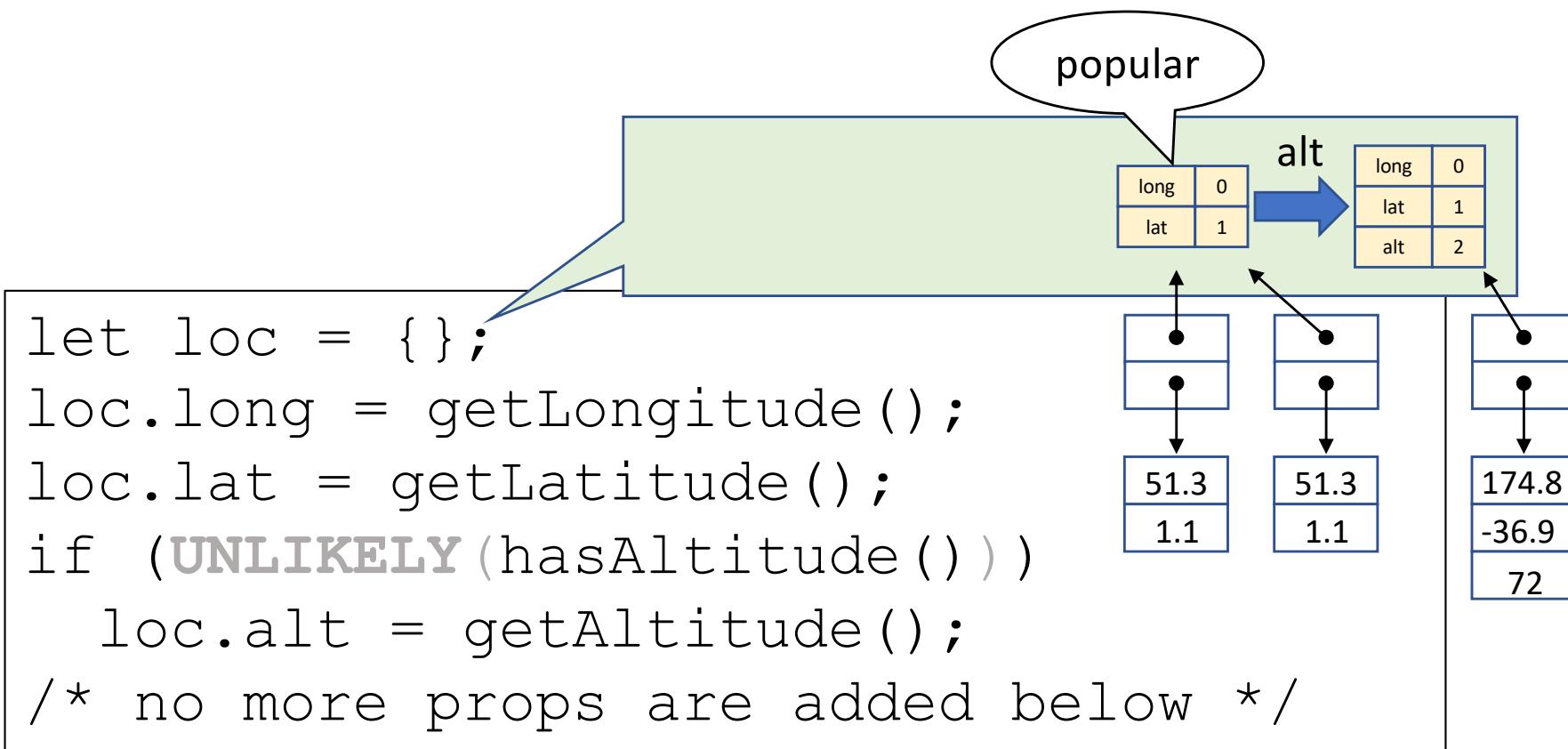
$K = 10$



# Optimisation 1': preserve popular HCs

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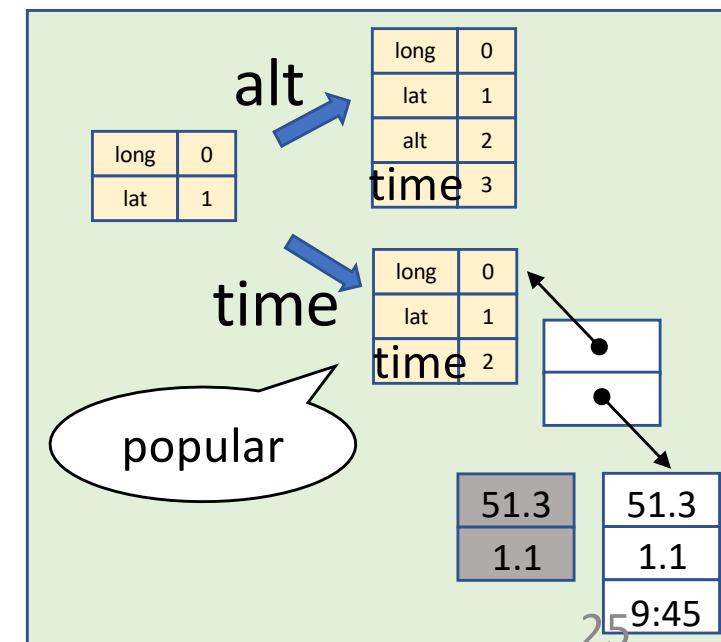
$K = 10$



# Motivating example for optimisation 2

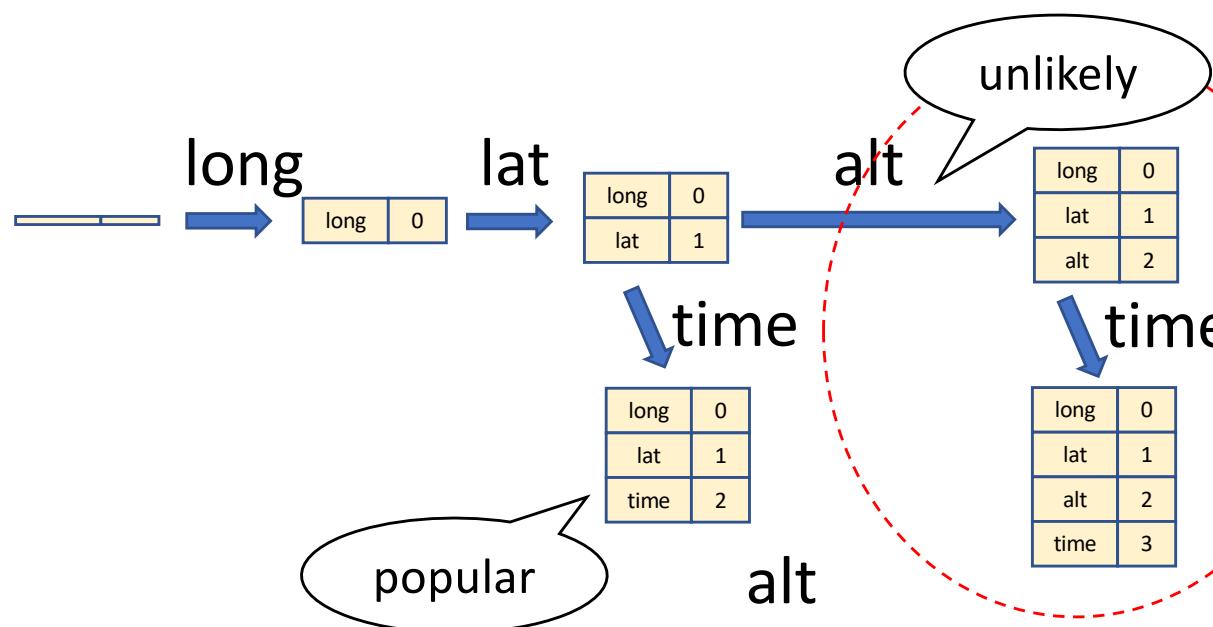
- Common properties are added after a branch
- Every object experiences property array re-allocation

```
let loc = {};  
loc.long = getLongitude();  
loc.lat = getLatitude();  
if (UNLIKELY(hasAltitude()) )  
    loc.alt = getAltitude();  
loc.time = currentTime();
```



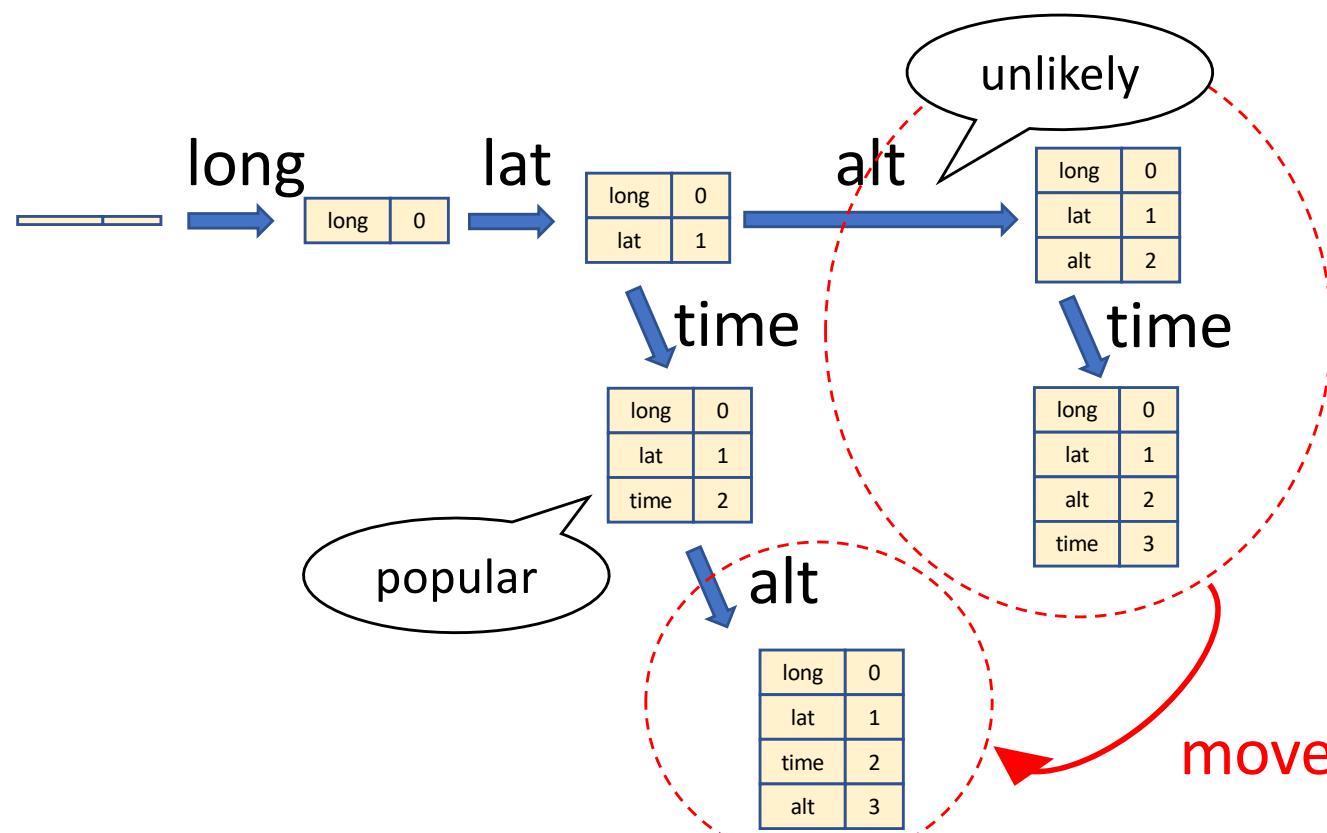
# Optimisation 2: moving branches

- Move “unlikely” branch before optimization 1



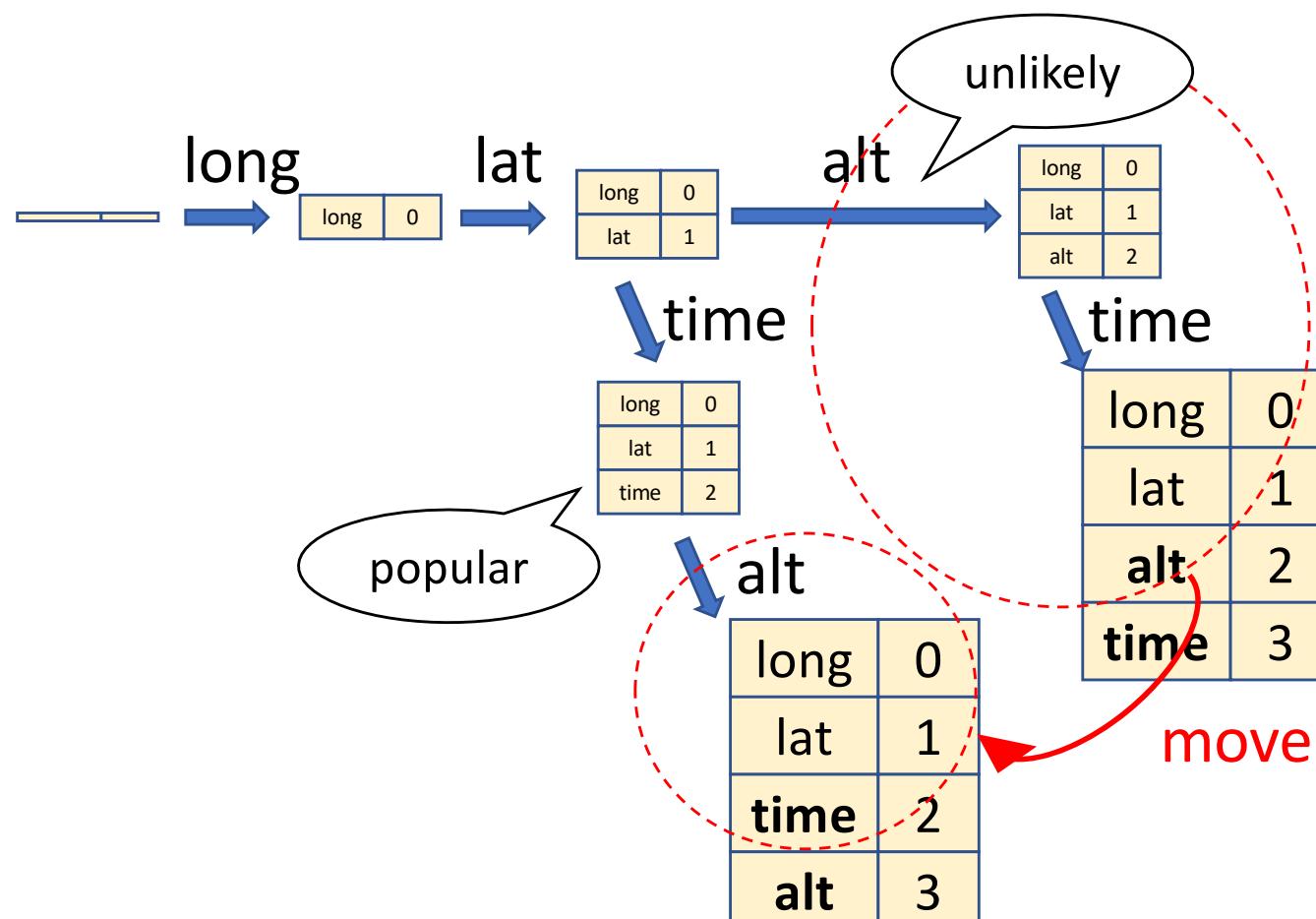
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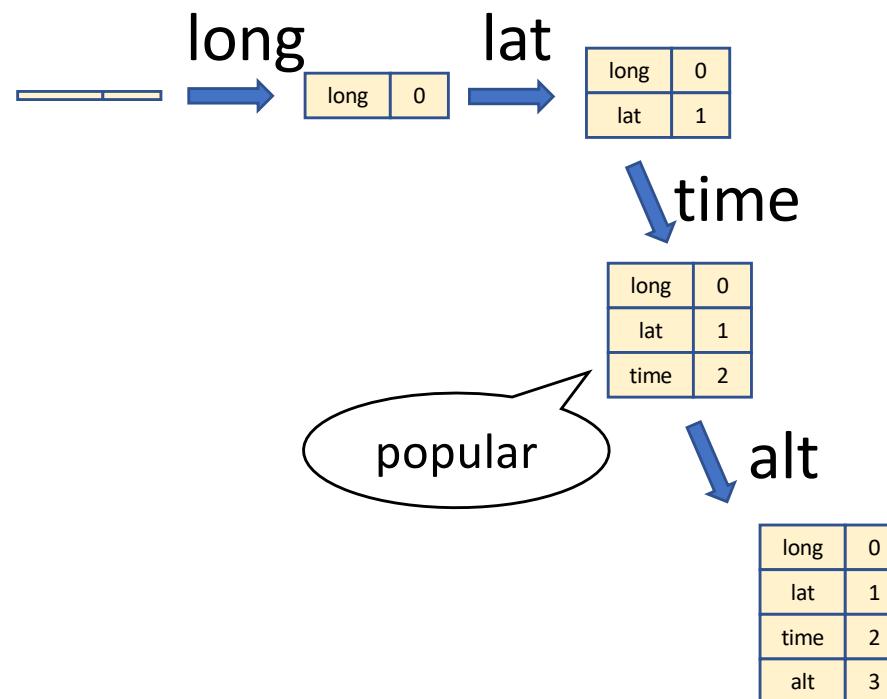
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# Optimisation 2: moving branches

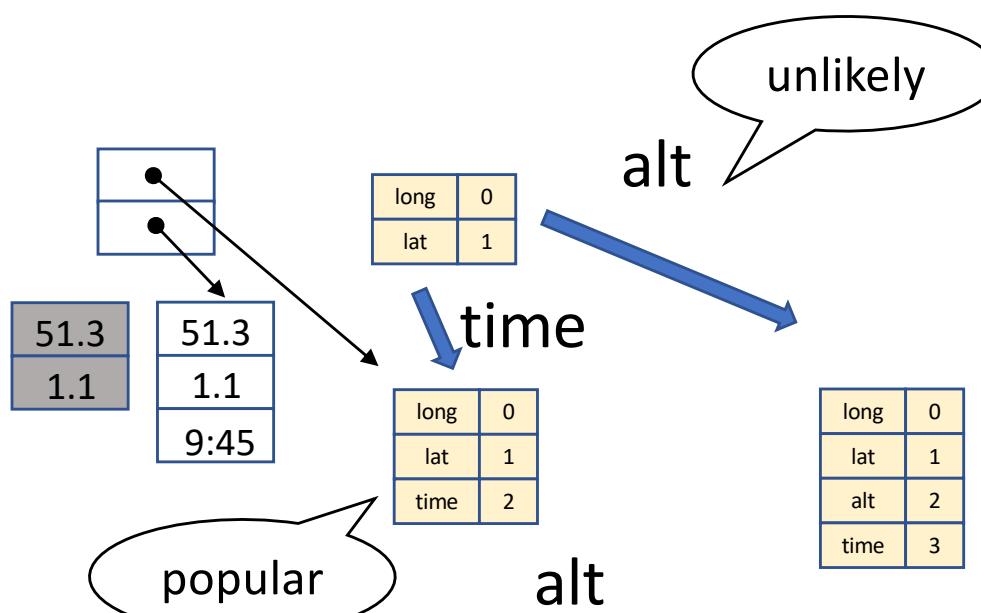
- Move “unlikely” branch before optimization 1
  - Linearise HCG



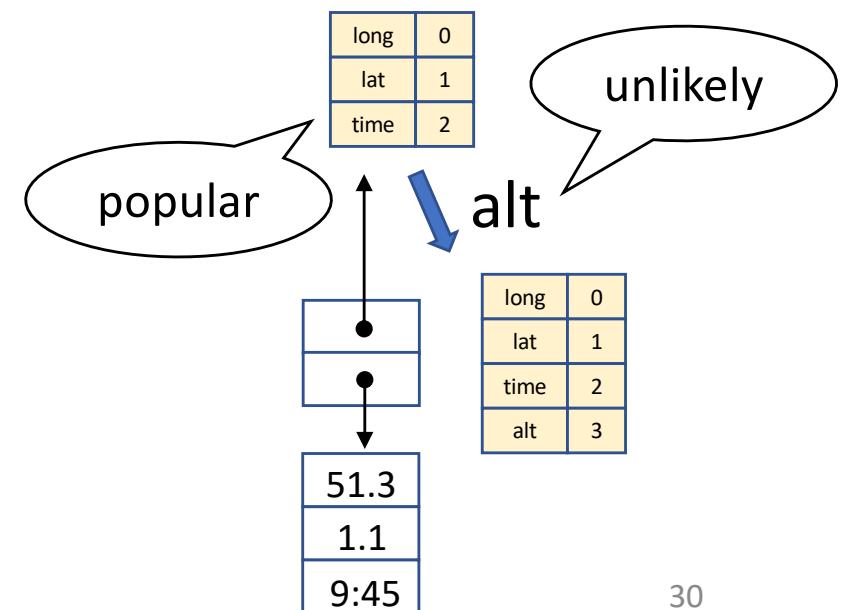
# Benefits of moving branches

- Encourage elimination of intermediate HCs
- Majority of objects are created with final layout

without moving branches



with moving branches

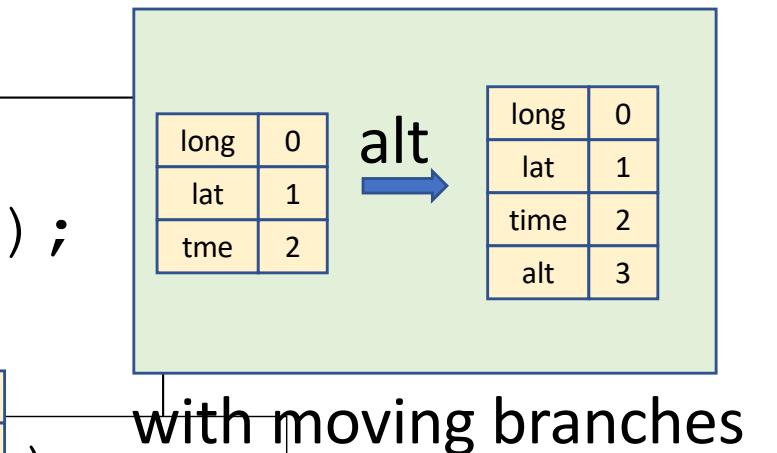
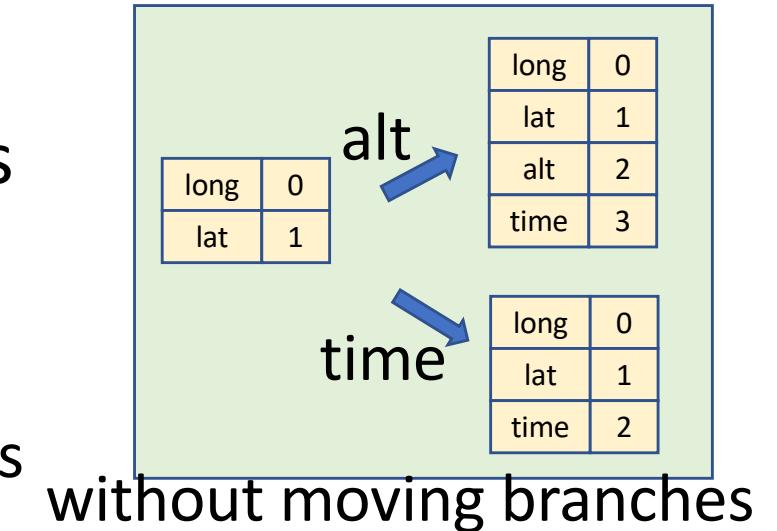


# Benefits of moving branches to inline cache

- Moving branch reduces variations of HCs
- Improves inline cache hit ratio
  - inline cache gives index if object has the same HC as cached

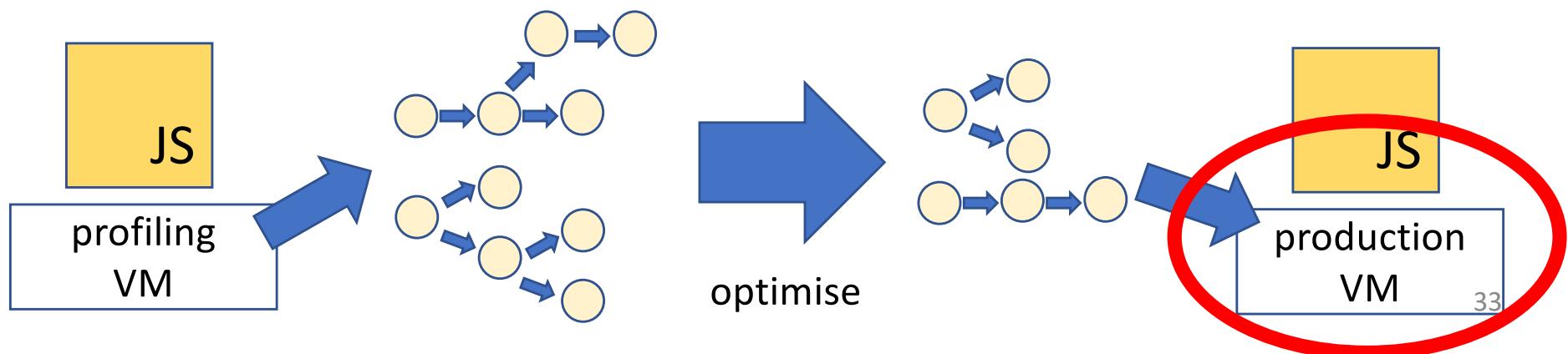
```
localTime(loc) {  
    tdiff = floor(loc.long / 15);  
    return loc.time + tdiff;  
}
```

```
if (obj->HC == time 2)  
    return obj->props [2];  
else  
    slow_path();
```



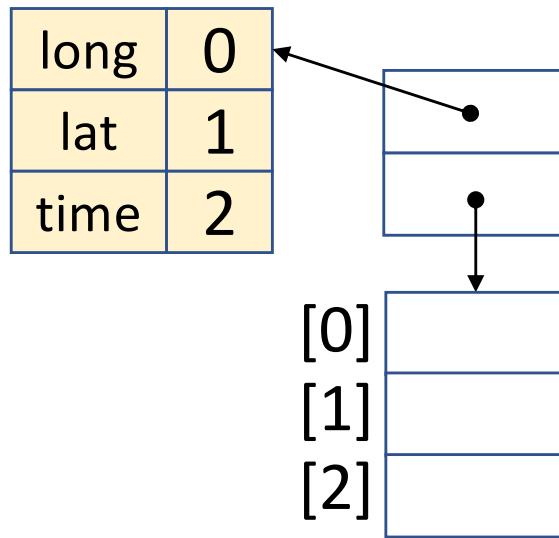
# Run-time optimisation

- Run-time optimisation relying on assumption that HCG is stable
  - in-object allocation
  - baking HCG into flash memory (future work)

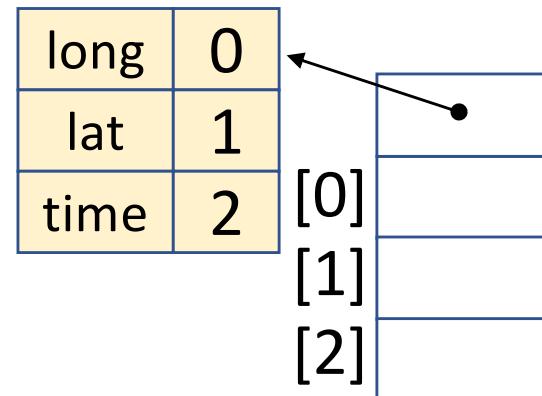


# in-object allocation

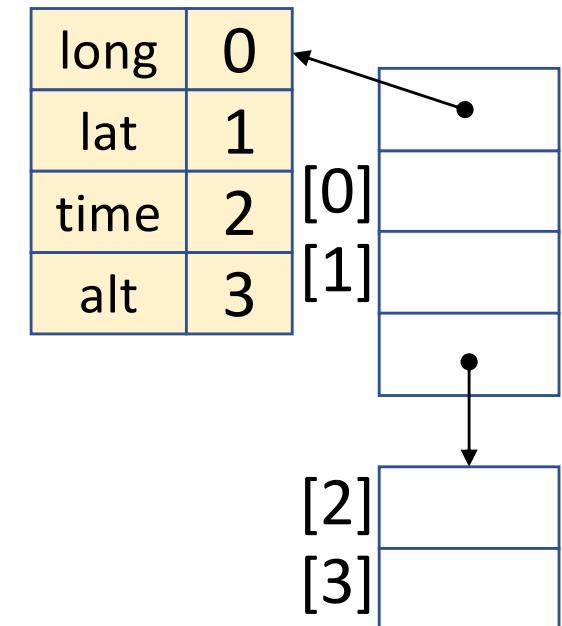
- Allocate all properties *in* object
  - Save space for indirect pointer
- In case of overflow,  
convert the last property area to indirect pointer



original version



in-object allocation

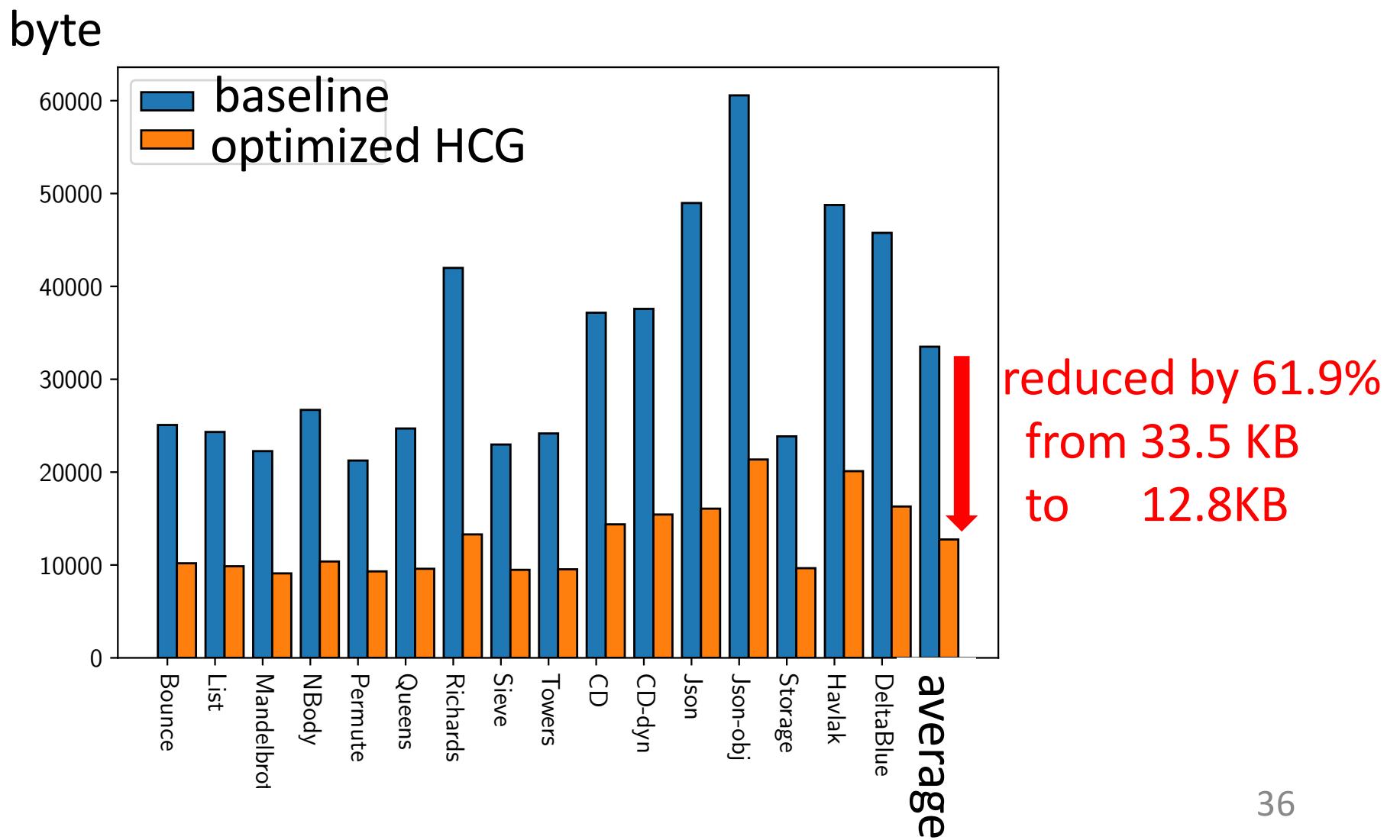


overflow

# Evaluation

- Implemented in eJSVM
- Are we fast yet benchmarks
  - original benchmarks
  - JSON-obj: uses an AST node object as a dictionary
  - CD-dyn: do not initialise future properties with NULL

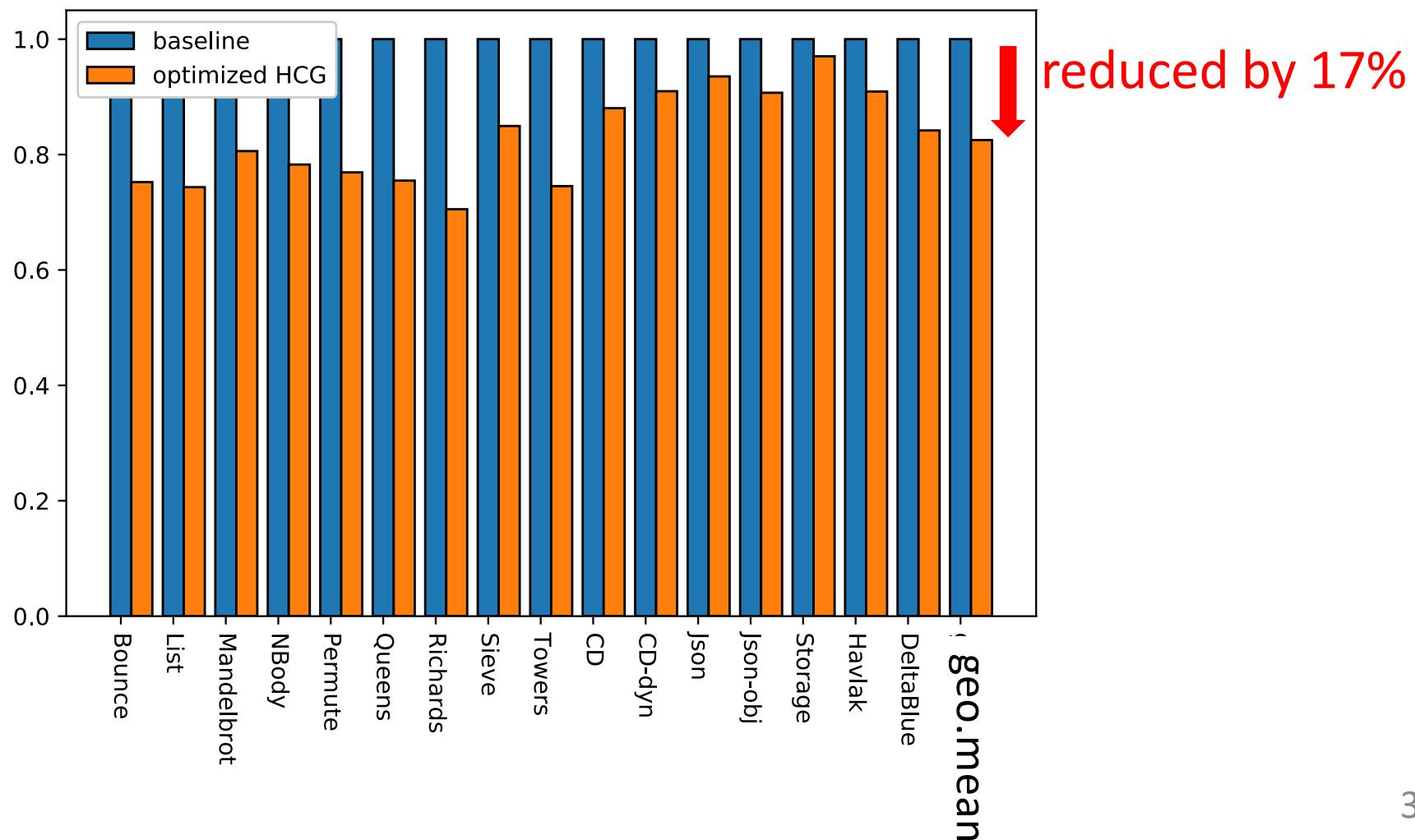
# Size of HC-related data



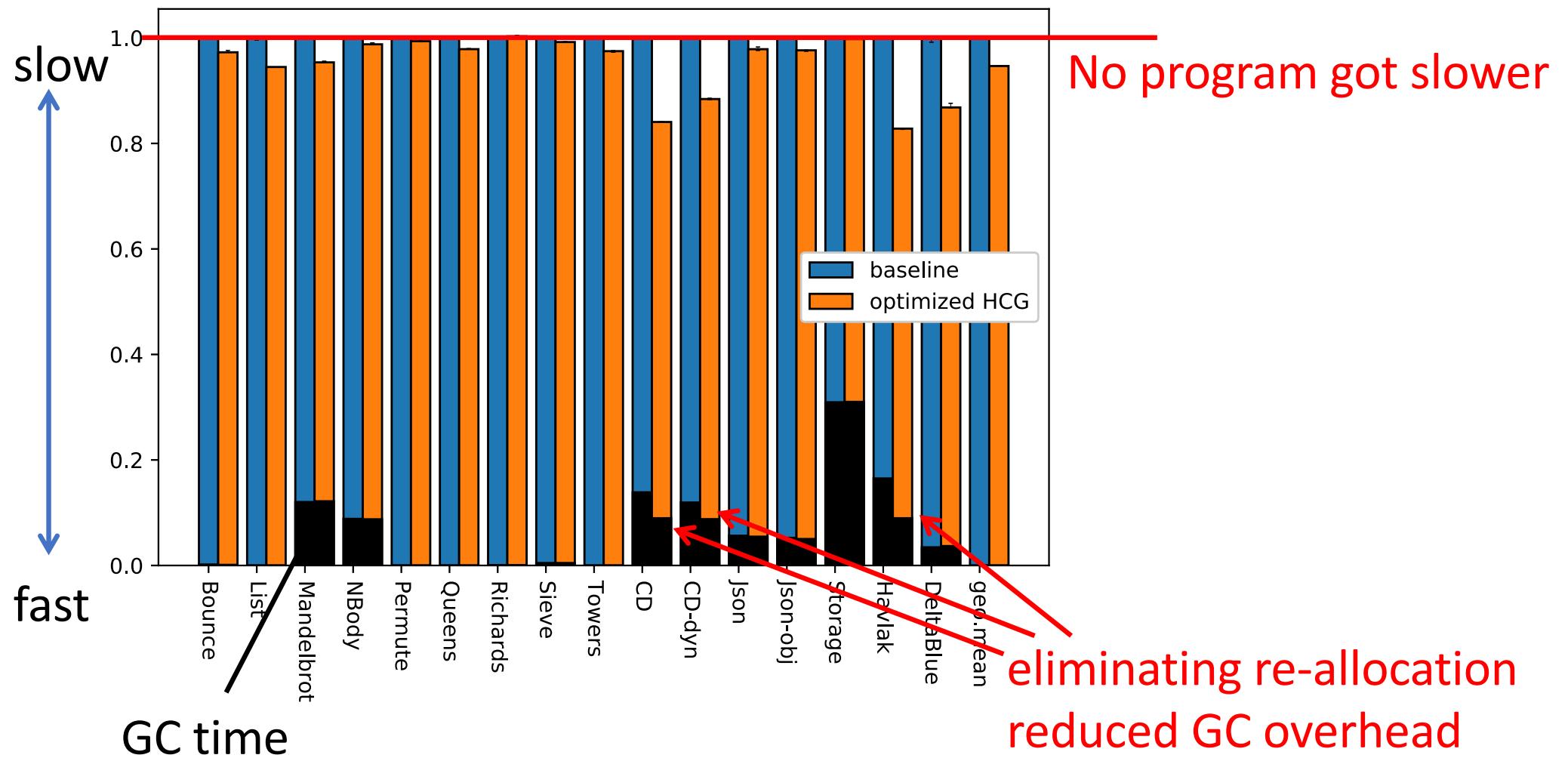
nomadised

# Maximum volume of all objects

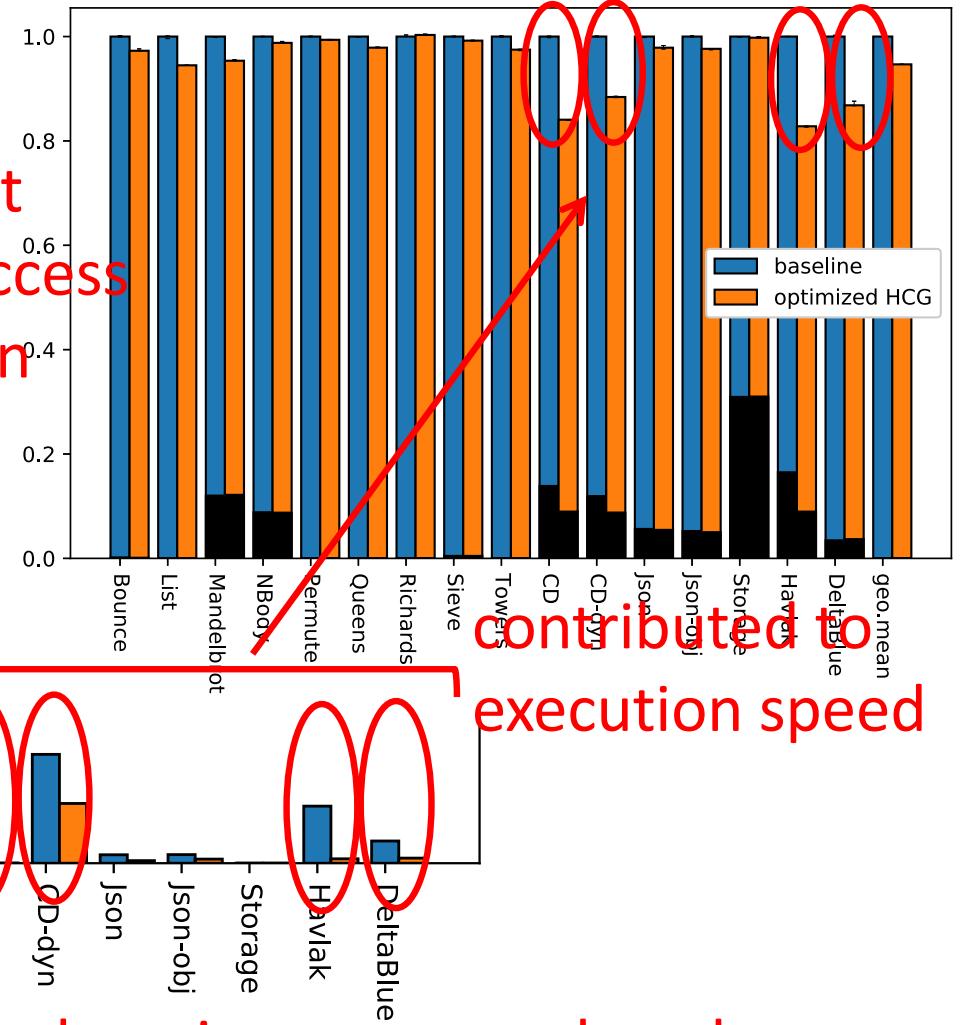
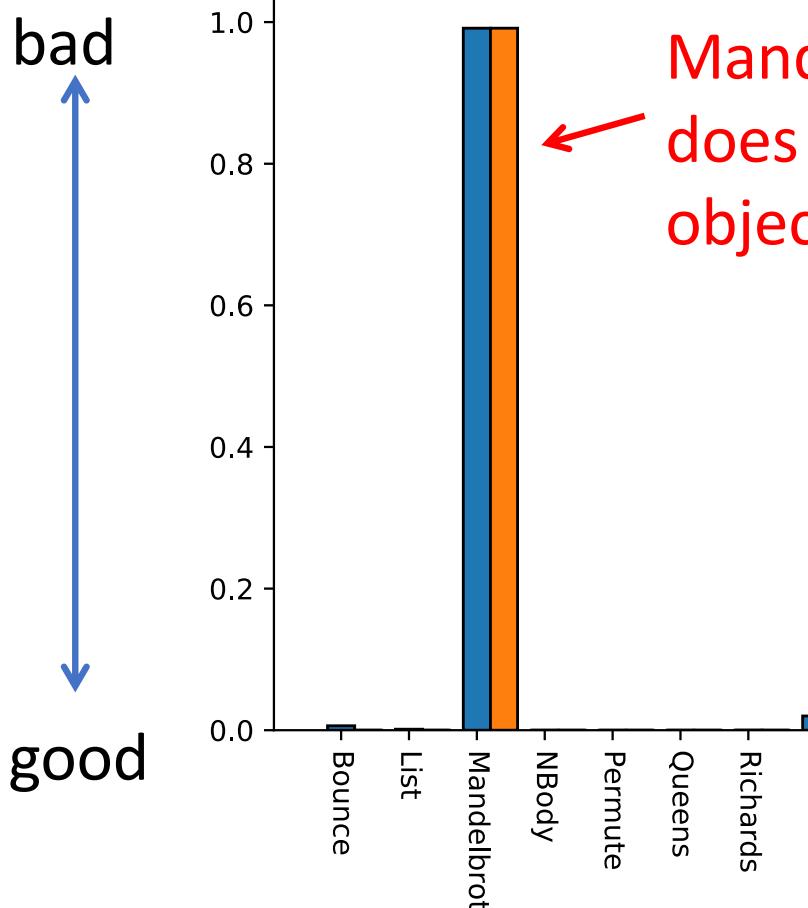
including HC-related data



# Normalised elapsed time



# Inline cache miss ratio



## cache misses were reduced

# Conclusion

- We proposed offline optimisation of HCG
  - Move “unlikely” branches
  - Eliminate intermediate HCs
    - preserve popular HCs
- Reduced HC-related data by 61.9% and footprint by 17%
- No program got slower
- Moving branches improved inline cache hit ratio